"Use of ATTACHed Interrupts in the UCSD p-System"

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ABSTRACT

Version IV.0 of the UCSD p-System supports concurrent processes with the use of the SIGNAL and WAIT primitive operations on semaphores. By use of the ATTACH primitive, semaphores may also be signalled by asynchronous interrupts.

This allows for application programs written in Pascal to respond to real-time external events. A sample program which does terminal emulation with file transmit and receive functions will be presented, along with a program which illustrates control of hardware functions of bit-oriented devices from a high-level language.

The UCSD p-System provides a convenient basis for software development on a large variety of microcomputers. Highly interactive, user-friendly programs may be constructed to fill a variety of single-user, personal workstation computing needs. However, many useful applications of microcomputers involve response to real-time events, and asynchronous control of devices which do not fall into the usual categories of printers, consoles, or character-oriented devices. The control of bit-oriented and analog devices by programs running under the UCSD p-System has been difficult or impossible until the release of Version IV.0.

Even without entering the realm of real-time machine control, a very common problem facing the programmer of a p-System personal workstation has been the construction of a 'terminal emulator' program, to allow communications to a remote host computer through a modern or other serial communications link. Appendix 1 contains a listing of a UCSD Pascal program which solves this problem for the specific case of communicating with Telemail, an electronic mail service available on GTE Telenet. The program is not represented as being completely general-purpose and userconfigurable, nor is it transportable to other p-System versions in object code form. However, modifications to the source text to fit other machines would be minor, provided the target implementation supports the attachment of external events to semaphore data structures.

For tutorial purposes, a much smaller and simplified version of the program is listed in Appendix 2. The basic algorithm is as follows:

repeat

{if a key is struck on the keyboard, send the character out to the modem}

{if a character is received from the modem, display the character on the screen} until finished;

A first attempt to code this in Pascal might look something like this:

repeat

read(keyboard,ch); write(remote,ch); read(remote,ch); write(output,ch);

until finished;

. . .which is quickly found to be of little use in practical time-sharing applications, since only one character from the host can be displayed for every character from the keyboard.

A solution which works on Version II.1 uses the Unit_Status system intrinsic to find out how many characters are in the input queue for the console and remote units:

repeat

Unit_Status(1,digits,1); if digits[0] <> 0 then begin for i := 1 to digits[0] do begin read(keyboard,ch);

write(remote,ch)

end;

Unit_Status(7,digits,1);
if digits[0] <> 0
then

begin

write(output,ch)

end;

until finished;

While this solution is adequate as presented, the overhead of the polling technique, combined with the desire to have more functions supported, make it difficult to keep up with a continuous 300 baud data stream. Practical terminal emulators with features such as file transmission and printer hardcopy became very large and complex, resorting to obscure Pascal coding tricks and to assembly-language external procedures in an attempt to sustain a 300 or 1200 baud data rate.

Version IV.0, as currently implemented on all TI 990 computers which support interrupts from character-oriented p-System devices, allows a 'full-function' terminal emulator program to be written with comparative ease, since language constructs are available which match the problem to be solved. This makes the main body of the terminal emulator appear effectively as:

repeat

{Nothing} until finished;

. . .since all the work is done by processes which are activated in response to external events.

These language extensions include a new data type (SEMAPHORE), a new block type (PROCESS), and five new pre-defined global procedures (ATTACH, SEMINIT, START, SIGNAL and WAIT). All except ATTACH are intended to be machine-independent constructs, and can be used on systems which do not support external asynchronous events. The use of these language features is described in the Version IV.0 Users' Guide, however, since the ATTACH intrinsic is somewhat machine-dependent, a description of the implementation on TI 990 computers is given here.

Just as a real processor needs to finish (or at least, bring to some orderly suspension point) the current machine instruction before handling an external interrupt, so must the p-machine. In practice, all of the pmachine interrupts which were generated by the real machine during the interpretive execution of the prior p-machine instruction are recognized before the next p-machine instruction is fetched. P-machine interrupts are selectively enabled by ATTACHing a SEMA-PHORE to an interrupt level, and disabled by AT-TACHing the value NIL to an interrupt level. The priority of the p-machine interrupt is determined by the priority of the PROCESS which is WAITing on the SEMAPHORE. In the current 990 implementation, 32 p-machine interrupts are supported; 64 will be supported in the IV.1 release in an effort to standardize the use of ATTACH on all p-System processors.

It is important to note the difference between a pmachine interrupt and an interrupt generated by hard-. ware and handled by the real processor. For example, the 990/10 receives a hardware interrupt on level 5 every 1/120th of a second. The machine code which handles this interrupt maintains a 32-bit integer which represents the time of day, but only on the full 1-second intervals does it generate the p-machine interrupt level 16, and then only if level 16 has been enabled by having had a SEMAPHORE ATTACHed to it. As a consequence, the DS990 Model 1 can run exactly the same Pascal PROCESS for handling the p-machine level 16 interrupt, even though the real-time clock on the Model 1 generates hardware level 2 interrupt every 250 milliseconds. These differences are handled in the machine code which services the interrupt, so that the Pascal programmer sees a consistent p-machine interrupt structure.

Another advantage to this de-coupling of hardware and software is that several p-machine interrupts can be generated by the same hardware interrupt under different circumstances, for example, a 0.1 second and a 10 second p-machine interrupt. Or more than one hardware interrupt could signal the same p-machine interrupt. Obviously, every hardware interrupt must have some machine code in place to handle the interrupt, but the amount of processing required in machine code could be as simple as acknowledgement of the interrupt bit and a generation of the appropriate p-machine interrupt. External (machine code) procedures are easily written which can be called from a Pascal program to effect CRU I/O and TILINE I/O.

Appendix 3 describes the requirements of a software driver which takes the output of a text formatter program and controls a letter-quality printer. In this case, the printer has a 24-bit parallel interface. The encoding of the print-stream information was designed with an assembly-language driver program in mind, but the Pascal solution given was much easier to write, and could be transported with a minimum of trouble. Further, the same Pascal source can be easily modified to interface a variety of different output devices to the same print-stream protocol, including the H-P 7220A plotter which produced the overhead projection texts.

The preparation and production of this paper on a personal workstation computer running the UCSD p-System is a (somewhat needless) demonstration of the usefulness of the p-System in an office environment. Likewise, the fact that the entire p-System can be maintained and re-generated on the same machine is no surprise to system software developers who have used the p-System. However, the introduction of several new factors in Version IV.0 now makes possible the application of the p-System to a large subset of the problems which previously required assembly-language operating systems for real-time process control. Use of ATTACHed Pascal PROCESSes, coupled with machine-code interrupt handlers, external I/O primitive routines, and the use of the Native Code Generator to produce machine code from Pascal p-code, will allow many real-time applications to be easily implemented in a high-level, transportable language and Operating System.

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Terminal emulator program for use with GTE Telenet and the Telemail service,
Uses interrupts from the keyboard, remote, and 1-second interval timer, ATTACHed to SEMAPHOREs which are WAITed on by several PROCESSes.
Fritz Whittington, Texas Instruments, Inc.
program telemail;
const
clocksig = 16;
keysig = $17;$
remsig = 18;
ar
ptr: integer; tick, godump,
waiting,

keyready, remready, eolecho, reply: semaphore; pid: processid; p: packed record case integer of a: packed array [0..80] of char); s: string[80]); 1: (2: (3: (4: (k: integer); x: integer; online: boolean); end; keych. remch. ret: packed array [0..1] of char; ch: char; dumping, quitting: boolean; i, j: integer; fname, mail, username, password. altuser, altpass, termtype, process busy: The OS does not tolerate a user program The OS does not tolerate a user program that has every process totally blocked. This process has lower priority than the main program, and cannot be blocked itself, since it doesn't wait on anything. If it is scheduled, it will be pre-empted sooner or later by another task of higher priority. begin while not quitting do {waste time}; signal(reply); {busy}; end process clocker; This process has the highest priority, and ensures that there will be an opportunity for task switching at one-second intervals. The signal(reply) ensures that the program doesn't become deadlocked due to a lost intervals. character echo. begin while not quitting do begin wait(tick); signal(reply); end; attach(nil,clocksig); end {clocker}; process dumper; This process lies dormant until a file is opened for transmission by procedure dof5. It then transmits the text a line at a time, putting a local copy on the console, and waits for the echo of the carriage return. (This delay avoids overrunning the Telenet input buffers, which were set up with human typists in mind.) A possible deadlock could occur if the echo is lost, but can be cured by entering a carriage return from the keyboard. During this process, the getrem process does not put echoed chars to the screen or log file. This process lies dormant until a file is

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begin

States:

s: string[255];

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var

repeat wait(godump); seminit(eolecho,0); if dumping then begin while not eof(gg) do begin readin(gg,s); writeln(s); {\$R-} {\$R+} unitwrite(8, ret, 1, , 16396); wait(eolecho); end; close(gg,lock); writeln('~~~Closing File: ',fname, '~~~'); end; writeln; dumping := false; end; until quitting; end {dumper}; procedure dof1; begin unitwrite(8,mail[2],(length(mail) - 1),,16396); unitwrite(8,ret,1,,16396); d {dofl}; end procedure dof2: begin {dof2}; end procedure dof3; begin {dof3}; end procedure dof4; begin end {dof4}; procedure dof5; begin attach(nil,keysig); We need to temporarily de-attach the keyboard signal, in order to do an ordinary readln of the file name. [_____ writeln; j = 99; {\$1-} repeat then begin fname := concat(fname,'.text'); reset(gg,fname); j := loresult; if j <> 0 then writeln('No such file'); close(gg); end else j := 0; until j = 0;{\$1+} if fname[1] <> chr(27) then begin jin
reset(gg,fname);
writein('~~~Transmitting File: ',fname,
'~~~'); dumping := true; signal(godump); end; ------Now, we can re-attach the semaphore to the keyboard interrupt so that the getkey process can handle characters.

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attach(keyready,keysig); This process is activated by a character being placed in the remote queue. The echoed chars lare normally placed in the logfile, except for control chars, and during file transmission, lsince the echo is unreliable from Telenet, and lanother copy of the text is not needed. end {dof5}; procedure dof6; begin unitwrite(8,altuser[2],(length(altuser)-1),,16396); unitwrite(8,ret,1,,16396); for j := 1 to 2000 do {waste time}; unitwrite(8,altpass[2],(length(altpass)-1),,16396); var unitwrite(8, ret, 1,, 16396); iord: integer; end {dof6}; begin procedure dof7; repeat wait(remready); begin if not quitting. Degin unitwrite(8, username[2], (length(username)-1),,16396); unitwrite(8, ret,1,,16396); for j := 1 to 2000 do {waste time}; unitwrite(8, password[2], (length(password)-1),,16396); unitwrite(8, ret,1,,16396); end [dof7]; then begin {get from remin} unitread(7,remch,1,,16384); ch := remch[0]; iord := ord(ch); if iord <> 10 then begin if not dumping then procedure dof8: begin {put to screen} begin unitwrite(8,logoff[2],(length(logoff)-1),,16396); unitwrite(8,ret,1,,16396); unitwrite(1,remch,1);
signal(reply); quitting := true; { dof8}; end; end end; if dumping and (iord = 13) then signal(eolecho); process getkey; if not dumping then This process is activated by a character being placed in the keyboard queue. The specialkey function returns TRUE if the key is one of the I'Fl' thru 'F8' keys on the 911 VDT, and trans-lates the keys to ASCII '1' thru '8'. if (iord = 13) or ((iord > 31) and (iord < 127)) then begin buff[ptr] := ch; ptr := ptr + l; if ptr = 1024 then begin ptr := blockwrite(ff, function specialkey: boolean; buff,2); ptr := 0; begin {specialkey} fillchar(buff,1024,chr(0)); specialkey := false; if ord(ch) in [146..153] then end: end: begin end ch := chr(ord(ch) - 97); else {quitting} begin specialkey := true; end; attach(nil,remsig); end; until quitting; end {specialkey}; end {getrem}; begin {getkey} repeat procedure initialize; wait(keyready); if not quitting then Initializes all the strings, and opens the next available logfile on the default volume. begin {get from keyboard} Unitclears the remote and console units. unitread(2,keych,1,,12); ch := keych[0]; if specialkey gin
 ret[0] := chr(13);
{The four strings which follow need to be changed
 to the appropriate values for each user. Note
 the unused ':' at the beginning, which is used to
 force word alignment for those PMEs that are
 sensitive to word alignment on unitwrites.}
 username := ':xxxxxxx';
 password := 'ixxxxxxx';
 altuser := ':xxxxxxx';
 logoff := 'ibye';
 mail := ':Mail';
 termtype := ':Dl';
 dumping := false;
 quitting := false; begin then begin case ch of 'l': dofl; '2': dof2; '3': dof3; '4': dof4; '5': dof5; '6': '7': dof6 dof7 1811 dof8: end {CASE} end else {not specialkey} begin {put to remote} [\$1-] unitwrite(8,keych,1,,16396); i := 0; := 0; end end repeat i := i + 1; p.s := '.x'; p.s[2] := chr(i + 64); fname := concat('termlog',p.s,'.text'); reset(ff,fname); else [if quitting] begin attach(nil,keysig); end; until quitting; d {getkey}; end j := ioresult; if j = 0 then process getrem; 30

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close(ff,lock)
                  else
           close(ff);
until j <> 0;
[$1+]
            rewrite(ff,fname);
fillchar(buff,1024,chr(0));
            ptr := blockwrite(ff,buff,2);
ptr := 0;
unitclear(1);
            unitclear(7);
            unitclear(8);
      end
                {initialize};
begin {main}
      initialize:
      write(chr(12));
writeln('Waiting for connection...');
writeln('214\748-0127 (300)...214\748-6371 (1200)');
      repeat
            p.online := false;
unitstatus(7,p.a,l);
      until p.online;
      unitclear(7);
      unitclear(/);
writeln('Use F1 for ''Mail'' command');
writeln('Use F5 to select transmit file');
writeln('Use F6 for Alternate user');
writeln('Use F7 for Username/password');
writeln('Use F8 for Quitting');
userclear(');
      unitclear(1);
     seminit(remready,0);
seminit(keyready,0);
seminit( godump,0);
seminit( reply,0);
seminit( tick,0);
      attach( tick,clocksig);
attach(remready,remsig);
attach(keyready,keysig);
      start( getkey,pid,800,142);
start( getrem,pid,500,140);
start( dumper,pid,500,129);
start(clocker,pid,500,145);
                         busy,pid,200,120);
       start(
       { Get Telenet's attention by sending 2 carriage returns at human speed, then define the video terminal}
      for i := 1 to 1000 do {waste time};
unitwrite(8,ret,1,,16396);
for i := 1 to 1000 do {waste time};
unitwrite(8,ret,1,,16396);
       for i := 1 to 1000 do {waste time};
unitwrite(8,termtype[2],2,,16396);
unitwrite(8,ret,1,,16396);
       repeat
       wait(reply);
until quitting;
       if ptr <> 0 then
    ptr := blockwrite(ff,buff,2);
close(ff,lock);
       signal(keyready);
signal(remready);
       signal(godump);
       attach(nil,clocksig);
attach(nil,keysig);
attach(nil,remsig);
       for i := 1 to 3000 do {waste time};
       writeln;
writeln('Returning to PascalSystem');
 end {telemail}.
```

```
Terminal emulator program tutorial.
Provides basic TTY emulation only.
     Uses interrupts from the keyboard, remote,
and 1-second interval timer, ATTACHed to
SEMAPHOREs which are WAITed on by several
                                                              remote,
     PROCESSes.
  Fritz Whittington, Texas Instruments, Inc.
program tutorial;
const
      clocksig = 16;
keysig = 17;
remsig = 18;
var
               tick,
        keyready,
remready: semaphore;
               pid: processid;
             keych,
             remch,
ch: char;
        quitting: boolean;
process busy:
   The OS does not tolerate a user program
that has every process totally blocked.
This process has lower priority than the
main program, and cannot be blocked itself,
since it doesn't wait on anything. If it
is scheduled, it will be pre-empted sooner
or later by another task of higher priority.
     begin
          while not quitting do {waste time};
     end
              {busy};
process clocker:
    This process has the highest priority, and
ensures that there will be an opportunity
for task switching at one-second intervals.
1
     begin
          while not quitting do
wait(tick);
{clocker};
     end
 process getkey;
 {-----
 This process is activated by a character being placed in the keyboard queue.
                              .
      begin {getkey}
          repeat
wait(keyready);
               {get from keyboard}
unitread(2,keych,1,,12);
               {put to remote}
unitwrite(8,keych,1,,16396);
           until quitting;
[ [getkey];
      end
 process getrem;
 This process is activated by a character being
 placed in the remote queue.
      begin
           repeat
                wait(remready);
               {get from remin}
unitread(7,remch,1,,16384);
{put to screen}
unitwrite(1,remch,1);
}
           until quitting;
```

Appendix 2

```
procedure initialize;
```

{getrem};

end

K H

```
begin
   quitting := false;
   unitclear(1);
   unitclear(7);
   unitclear(8);
end {initialize};
```

Unitclears the remote and console units.

begin {main}

```
initialize;
```

```
seminit(remready,0);
seminit(keyready,0);
seminit( tick,0);
attach( tick,clocksig);
attach(remready,remsig);
attach(keyready,keysig);
start( getkey,pid,800,142);
start( getrem,pid,500,140);
start(clocker,pid,500,145);
start( busy,pid,200,120);
```

repeat wait(tick); until quitting;

{Note that there is no provision for quitting, and that the (higher priority) process clocker is also waiting on tick.}

end (tutorial).

T2WPUNIT Interface Description

The text formatter program generates a print-stream which may be directed either to a file or to a printer. If the print-stream is directed to a file, it may subsequently be printed with the de-spooling program. Both of these programs interface with the physical printer by means of a unit T2WPUNIT which contains a signal public procedure WPINTERFACE(ch: char) to which the print-stream is sent byte-by-byte.

```
UNIT T2WPUNIT;
INTERFACE
      procedure wpinterface(ch: char);
IMPLEMENTATION
      var
             first_time: boolean;
      procedure wpinterface;
      begin
      if first_time then
             begin
             first_time := false;
             {here initialize printer}
             end:
      {here process print-stream byte ch}
      end {wpinterface};
BEGIN (T2WPUNIT)
first_time := true;
***
END [T2WPUNIT].
```

For documents printed in fixed-pitch mode, the print-stream is a simple ASCII stream. Each line of print is sent from left to right and terminated by a CR. The CR is usually followed by one of more LF's according to the amount of paper movement required between lines. However, where overprinting is required (e.g. for underlining), the CR is not followed by a LF. For fixed-pitch printers which are not capable of overprinting (because they have no non-advancing end-ofline function), the WPINTERFACE driver can throw away all characters after a CR until a LF is encountered. When single-forms mode is selected, the printstream contains an ASCII group-separator character GS to signify that the driver should pause for a new sheet of paper to be positioned (or possibly should activate an automatic sheet feeder).

For documents printed in variable-pitch mode, the print stream contains tightly encoded information on horizontal printhead movement, vertical paper movement, print-element character selection, hammer energy, special print effects, and single-forms control. Such variable-pitch information is always introduced in the print-stream by the ASCII sequence NUL-NUL-SO to indicate to the WPRINTERFACE driver that variable-pitch mode is required. The driver must decode the print-stream and produce the requested printer actions. The following page defines the contents of the print-stream when in variable-pitch mode.

Appendix 3

Variable-Pitch Print-Stream Description

Drivers are instructed to enter variable-pitch mode by the sequence NUL-NUL-SO, and to leave variablepitch mode by the sequence NUL-NUL-SI. Once in variable-pitch mode, the driver must respond to the following character sequences:

Type 1: (2 bytes)	lrrr rrrr	eeem mmmm
Type 2: (2 bytes)	01dd mmm	mmm mmm
Type 3: (1 byte)	00cc cccc	

where:

rrrrrrr = 7-bit character (code or rotation)

eee = 3-bit hammer energy (zero if n/a)

m...m = 5-bit or 12-bit movement distance:

in 1/120-inch increments for horizontal movement in 1/48-inch increments for vertical movement dd = 2-bit movement direction:

00 = right (forward tabulation) 01 = left (reverse tabulation) 10 = down (forward line feed) 11 = up (reverse line feed)

ccccc = 5-bit special action code:

NUL = no action GS = pause after single form BS = set direction = backward HT = set direction = forward CR = home printhead, set direction = forward ESC = next Type-1 sequence defines underline SI = return to fixed-pitch mode

Any Type-3 byte value from 30 to 3F hexadecimal sets special-effects:

byte = 0011 xdbu where:

u = underscore mode b = boldface mode

- d = double-strike mode
- x = (not defined yet)

Type-1 sequences cause the print-head to move mmmmm 120ths of an inch to the left or right (whichever was set by the last Type-2 or Type-3 sequence), and then the character rrrrrrr to be printed with hammer energy eee. {SN+}
UNIT T2WPUNIT;
(UNIT for the NEC Spinwriter with parallel card
and using ABSOLUTE SPOKE mode)
INTERFACE
PDCCEDURE wpinterface(chi char); IMPLEMENTATION const wpcru = 64; [base address of interface card] [cru output bits] restore = 12; {sbo to restore, sbz to run} select: = 16; [sbo to select] ribift = 20; [sbo for lower ribbon part [b] halfsp = 11; [lsb of horizontal movement] waybit = 10; [sbz = right or down (head wrt lpback = 13; [sbo to read bits back for tes pwstb = 18; [print wheel strobe] ofstb = 17; [paper feed strobe] {sbo to restore, sbz to run}
{sbo to select}
{sbo for lower ribbon part (black)}
{lsb of horizontal movement}
{sbz = right or down (head wrt paper)}
{sbo to read bits back for test}
{print wheel strobe}
{carriage strobe} riblit = 20; halfsp = 11; waybit = 10; lpback = 13; pwstb = 18; pfstb = 17; carstb = 19; (cru input bits) papout = 22; ribout = 21; pcheck = 20; pready = 16; pfready = 17; pwready = 18; carready = 19; pe (true if paper out)
(true if ribbon out)
(true if printer in check)
(true if printer ready)
(true if paper feed ready)
(true if print wheel ready)
(true if carriage ready) 19; [true if carriage ready; = 0..1; = 0..3; = 0..7; = 0..7; = 0..12; = 0..63; = 0..4095; = 0..4095; = 0..4095; = 0..4095; = (right,left,down,up); [direction of head wrt paper] = (action,s_effects); = (action,s_effects); = (prefix,ud],ud2,ud3,uv4,ud5,ud6,ud7, setback,setnorm,ud10,ud11,ud12,home,entervp,exitvp, ud16,ud17,ud13,ud19,ud20,ud21,ud22,ud23, ud26,ud25,ud26,setus,ud28,hold1t,ud30,ud31); = packed record [for byte-swapping] r_byter ubyte; end; = nacked record [for TYPE 1 commands] type onebit twobits threebits nibble fivebits sixbits sevenbits ubyte twelvebits bits typekind waykind t3kind actionkind word0 {for TYPE 1 commands} fivebits; threebits; sevenbits; boolean; packed record wordl movement: energy: rotation: istypel: end; packed record word2 {for TYPE 2 commands} twelvebits; movement: direction: whichtype: waykind; typekind; end; packed record word3 (for TYPE 3 commands- actions) ubyte; actionkind; t3kind; resl: t3action: whicht3: whichtype: typekind; end; packed record {for TYPE 3 commands- effects} word4 {word5 is just an integer and is denoted by (2), word6 = packed record {for 1355 Diablo in Rib. Opt. 2} spoker sevenbits; ribbon: threebits; energy: twoblts; resit nibble; end; {word7 is a packed array of bits} word8 = packed record {for NEC Spinwriter in Absolute mode} abs: boold energy: three resl: five: abs: boold energy: three end; case integer of 0: (w0: word0); 1: (w1: word1); 2: (w2: word2); 3: (w3: word3); 4: (w4: word4); 5: (15: integer); 6: (w6: word6); 7: (bit: bits); 8: (w8: word8); end; boolean; threebits; fivebits; urec endi

var ch: char; cacked: boolean; pat packed array [0..10] of integer; begin (check_ready) cpuidle; (allow cpu to idle until interrupt) if ((not testbit(pready)) or (testbit(papout)) or (testbit(ribout)) or (testbit(pcheck)) then then begin unitclear(l); acked := false; while not acked do
begin
write(chr(13+128)); {Return without auto linefeed}
write('Printer needs attention- press <space>',chr(7));
cpuidle;
pa[0] := 0;
unitstatus(1,pa,1); {get number of keys buffered for input}
if pa[0] <> 0 then read(keyboard,ch);
if ch = ' ' then acked := true;
end; if ch = ' ' then acked := true; end; unitclear(1); gotoxy(0,100); writeln('Correct printer condition (paper, ribbon, cover)'); writeln; writeln('Printing will resume when <enter> is pressed'); repeat read(keyboard,ch) until ch = chr(160); j; repeat read(ke end; end; {check_ready} procedure hardinit; var procedure hardinit; var i: integer; begin (hardinit) cruword.i5 := 0; loadcru(cruword.i5,16); [clear all data line setbit(select); {select the printer] setbit(restore); {tell it to restore} clearbit(pistb); {deactivate strobe l clearbit(carstb); {deactivate strobe l setbit(ribit); {raise the ribbon} for i := 1 to 500 do; {waste some time} clearbit(restore); {let go of restore l i t= 1; while (i < 3000) and (testbit(pready) = false) do i := i + 1; {clear all data lines} {select the printer} {tell it to restore} {deactivate strobe line} {deactivate strobe line} {deactivate strobe line} {raise the ribbon} {waste some time} {let go of restore line} do i := i + l; if not testbit(pready) then then begin writeln('Cannot open printer'); end; check_ready; end; [hardinit] procedure print_it(k: urec); var and (not testbit(pwready)) begin check_ready; timeout := timeout + 1; timeout := timeout + 1; end; setbit(pwstb); clearbit(pwstb); end; {print_it} procedure move_carriage(howfar: integer; whichway: waykind); var timeout,tempmove,bigmove,lsb,lead: integer; timeout,tempmove,bignove,..., begin [move_carriage] if whichway = right then horpos t= horpos + howfar else horpos t= horpos - howfar; check_ready; if not underlining then then begin while howfar > 0 do
begin
if howfar > 1023 then tempmove i= 1023 else tempmove i= howfar;
howfar := howfar - 1023;
bigmove := tempmove div 2;
lsb i= tempmove mod 2;
if lsb = 1 then setbit(halfsp) else clearbit(halfsp);
case whichway of
right: clearbit(waybit);
left: setbit(waybit);
end; do end; loadcru(bigmove,10); timeout := 1; while (timeout < 30000)</pre> and (not testbit(carready)) begin check_ready; timeout := timeout + l; endi end; setbit(carstb); clearbit(carstb); l; {while howfar > 0} end:

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{of then clause on "if not underlining" }
{if underlining} end eise begin while howfar > 0 de begin printit(usrec); case whichway of right: clearbit(waybit); left: setbit(waybit); end: end; bigmove := 0; {that is 0/60ths} setbit(halfsp); {+ 1/120th } howfar t= howfar - 1; loadcru(bigmove,10); **movit t= 1; timeout i= 1; while (timeout < 30000)</pre> and (not testbit(carready)) do begin check_ready; timeout := timeout + 1; end; setbit(carstb); setDit(carstb); clearbit(carstb); end; {while howfar > 0} end; {of the "if not underlining"} end; {of the "if not underlining"} end; {move_carriage} procedure move_paper(howfar: integer; whichway: waykind); var timeout *comments. var timeout,tempmove: integer; begin {move_paper} if whichway = down then vrtpos := vrtpos + howfar else vrtpos := vrtpos - howfar; check_ready; setbit(halfsp); case whichway of {DOWN means printhead down, down: clearbit(waybit); upt setbit(waybit); end; {DOWN means printhead down, paper UP} end: begin check_ready; timeout := timeout + 1; timeout i= ti end; setbit(pfstb); clearbit(pfstb); end; end; clearbit(pist); end; end; end; (move_paper} procedure home_head; begin (home_head) hdir := left; move_carriage(horpos,hdir); hdir := right; end; (home_head) procedure turn_vp; begin {turn_vp} fixedpitch := false; end; (turn_vp) procedure turn_fp; begin {turn_fp} begin {turn_fp} fixedpitch := true; end; (turn_fp] procedure initialize; begin itvrn_fp fixedpitch := true; end; (turn_fp] procedure initialize; begin intypel := false; intypel := false; fixedpitch := true; nulcount := 0; inusdef := false; setbase(wpcru); vdir := down; hdir := right; usrec.if := 0; usrec.wl.istypel := true; bold := false; dubstrike := false; Dold := false; dubstrike := false; underline t= false; vrtpos := 0; horpos t= 0; hardinit; turn fa. notypos 1= 0; hardinit; tur_fp; home_head; end; {initialize} procedure ex_pause; var cg: char; begin [ex_pause] writein('Printing suspended by [single forms]'); writein('Printing will resume when <enter> is pressed'); repeat read(keyboard,cg) until cg = chr(160); end; [ex_pause]; var saveusf: hoc/results var
saveusf: boolean;
bdir: waykind;
howfar: integer;
begin [ex_type]]
if inusdef then then
begin
 usrec.i5 := outrec.i5;
 inusdef := false;
 exit(ex_typel); end; howfar := outrec.wl.movement;

move_carriage(howfar,hdir);
print_it(outrec);
if dubstrike and (not bold) then print_it(outrec);
if bold then then begin case hdir of right: bdir := left; left: bdir := right; end; saveusf := underlining; underlining: false: saveusf := underlining underlining := false; move_carriage(1,hdir); print_it(outrec); print_it(outrec); move_carriage(1,bdir); print_it(outrec); underlining := saveusf; end; end; [ex_type1] procedure ex_type2; var howfare :=procedure _____ var howfar: integer; begin (ex_type2) case outrec.w2.direction of right: hdir := right; left: hdir := left; down: vdir := down; up: vdir t= up; up, val. 1 = up; end; if (outrec.w2.direction = down) or (outrec.w2.direction = up) then --o+n
howfar i= outrec.w2.movement;
move_paper(howfar,vdir);
end else else begin howfar := outrec.w2.movement; move_carriage(howfar,hdir); end; {ex_type2} procedure ex_type3; begin {ex_type3} if nulcount = 2 then begin nulcount := 0; end; end; if outrec.w4.whicht3 = action then begin case outrec.w3.t3action of se outrec.w3.t3action of
prefix: nulcount := nulcount + 1;
setback: hdir := left;
setnorm: hdir := right;
home: home.had;
entervp: turn_vp;
exitvp: turn_fp;
setus: inusdef := true;
holdit: ex_pause;
d; {case} end; end; {case; end else begin if outrec.w4.bf_on then bold := true else bold := false; if outrec.w4.ds_on then dubstrike := true else dubstrike := false; if outrec.w4.us_on then underlining := true else underlining := if outrec.w4.u faise; end; end; [ex_type3] procedure process_it; var var pat packed array [0..1] of char; begin [process_it] if fixedpitch and (ch <> chr(0)) and (nulcount <> 2) theo then then
 begin
 pa[0] := chr(ord(ch));
 unitwrite(6,pa,1,,12);
 exit(process_it);
end;
if intypel
then
beevin then
begin
outrec.w0.r_byte := ord(ch);
intypel := false;
ex_typel;
exit(process_it); end; if intype2 then then
begin
outrec.w0.r_byte := ord(ch);
intype2 := false;
ex_type2;
exit(process_it); end; inrec.w0.l_byte := ord(ch); if inrec.w2.whichtype = type3 then begin over outrec.w0.i_byte := ord(ch); ex_type3; _exit(process_it); end else begin if inrec.wl.istypel then begin outrec.w0.l_byte i= ord(ch); intypel i= true end else begin outrec.w0.l_byte := ord(ch); intype2 := true;

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	d;				x	r5	; Execute it	
end; end;	{process_i	it]			rtwp .proc	loadcru,2	; Done	
begin	{wpinter				.rei	crubase, ws		
	rst_time t			;pascal	declaratio	on;		
beg				; proc	edure load	cru(i,j: intege	er); external;	
	st_time t: tialize;	: false;		; whe	re 1 is the	e 16-bit patter	n desired, j is number of bits	
end				, 1990 in	struction	ls: li	rx,l	
proce	ss_it;			1			r rx,j	
end;	{wpinter	face)		;	mov		D	
BEGIN (T	ZWPUNII}	; {UNIT initial	i)		mov	*r10+,ws+2 *r10+,ws+4	; Pop # of bits into new ri ; Pop data word into new r2	
***;	•	, town initial	ization)		blwp	start	; Go do it	
(no UNIT	terminati	on }			Ь	*r11	; Return to pascal	
END T2W	PUNIT).			start inst	.word Idcr	ws,ep r2,0	Lost suction much	
, : Machin	e code for	the external f	unctions	ep	mov	crubase, r12	; Instruction mask ; Get current base from private storag	
;				- 1	THO V	inst,r5	; Puts 'ider r2,0' instruction in r5	
	.proc	machine	· · · · · · · · · · · · · · · · ·		andi	r1,000fh	; Mask # of bits to 0-15	
	limi	1 (38	; To simulate the effect of a ; Machine check (level 2 interrupt)		sia soc	r1,6	; Shift to spot required in instructio	
	blwp b	*rii	, maentile eneer (refer 2 interrupt)		x	r1,r5 r5	; Or into r5 ; Execute it	
	.proc	cpuidle			rtwp	-	Done	
	idle	• • •			.proc	storecru,2		
	Ь	*rii	·		•ref	crubase,ws		
	.proc .dei	setbase,l crubase,ws			declaratio		eger;): integer); external;	
; pascal	declaratio			-; whe	re i is the	e 16-bit result	area, j is number of bits	
; proce	dure setba	<pre>se(i: integer);</pre>	external;	;				
; wher	e i is the	crubase desire	d (decimal)	;990 in	struction		rx,j	
; ;990 ins	truction i	s: li	r12,i	i		mov	rx,i	
;				,	mov	*r10+,ws+2	; Pop # of bits into new rl	
	mov	*rl0+,crubase			mov	*r10+,ws+4	; Pop address of data into new r2	
crubase	b .word	*r11 0	; Return to pascal		blwp b	start	; Go do it	
WS	.block	32	: Workspace for the others	start	word	*rll ws,ep	; Return to pascal	
	.proc	setbit,1		inst	ster	*r2,0	; Instruction mask	
	.ref	crubase, ws		ep	mov	crubase,r12	; Get current base from private storag	
	declaratio	t(i: integer);	external		mov andi	inst,r5	; Puts 'stor *r2,0' instruction in r5	
; wher	e i is the	bitnumber desi	red		sla	rl,000fh rl,6	; Mask # of bits to 0-15 ; Shift to spot required in instructio	
;					50C	r1,r5	; Or into r5	
;990 ins	truction i	si sbo	i		x	r 5	; Execute it	
,	mov	*r10+,ws+2	; Pop stack into new rl		rtwp .iunc	testbit,1	; Done	
	blwp	start	; Do the stuff		.ref	crubase,ws		
	Ь	*r11	; Return to pascal	;pascal	declaratio	on;		
start .word ws,ep inst sbo 0 ; Instruction mask			; function testbit(i: integer): boolean; external;					
1		crubase, ri2	; Get current base from private storage	; where i is the bitnumber				
•	moy	inst, r5	; Op code for sbo	;990 in	struction i	is: tb	i •	
	andi	rl,00ffh	; Insure parm passed is 00 in first byte	;		(ret	urns true if bit is 1)	
	soc x	r1,r5 r5	; R5 now has proper sbo instruction ; Execute it	;	mov	# # 10 mm - 0		
	rtwp		; Done *		ai	*r10,ws+2 r10,2	; Pop # of bits into new rl ; Point to word provided on stack	
	proc	clearbit,1		; Note-		pointer is lef	t pointing to result word	
	.ref	crubase,ws			blwp	start	; Go do it	
	declaratio	on; ·bit(i: integer)	· external·		mov	ws,*r10	; Push r0 of new ws	
		bitnumber desi		start	b .word	*r11	; Return to pascal	
;				inst	tb	ws,ep 0	; Instruction mask	
;990 ins	truction i	sı sbz	i	ep	mov	crubase,r12	j Get current base from private storag	
	mov	*r10+,ws+2	; Pop stack into new rl		cir	rO	; Assume bit is false	
,		start	; Do the stuff		andi mov	rl,00ffh inst,r5	; Insure displacement is in range ; Get copy of instruction	
,	DIWP	*r11	; Return to pascal		SOC	r1,r5	; Or in displacement	
,	blwp b							
start	.word	ws,ep			x	r 5	; Execute test bit	
inst	b .word sbz	ws,ep 0	; Instruction mask		jne	\$1	; Jump if bit was false	
	b ,word sbz mov	ws,ep 0 crubase,r12	; Instruction mask ; Get current base from private storage		jne inc		; Jump if bit was false ; If bit was 1, r0 := 1	
inst	b .word sbz	ws,ep 0	; Instruction mask	\$1	jne	\$1	; Jump if bit was false	

The Two Faces of UCSD Pascal

By Rich Gleaves Volition Systems

Rich Gleaves of Volition Systems submits the slides from a talk of his. I (as usual) dropped out where and when the talk was. I have included the first few slides of his talk of which the outline was as follows:

THE TWO FACES OF UCSD PASCAL UCSD PASCAL _ INDUSTRY IMPACT UCSD PASCAL - HISTORY UCSD PASCAL SYSTEM VERSIONS UCSD PASCAL vs. STANDARD PASCAL UCSD PASCAL EXTENSIONS PROGRAM SEGMENTATION SEPARATE COMPILATION - UNITS UCSD I/O HIERARCHY INTERACTIVE I/O RANDOM ACCESS FILES UNIT I/O STRINGS BYTE ARR'AY MANIPULATION DYNAMIC STORAGE PROCEDURE TERMINATION EVEN MORE TRICKS RECORD AND ARRAY COMPARISON

If you would like to have the bodies for all these slides, please contact Volition Systems. ed.

THE TWO FACES OF UCSD PASCAL

• Friendly beginner's language

Used at UCSD to teach introductory computer programming to non-science students. UCSD Pascal's